# Logics for Specification

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### CSP & CASL

### Modelling Concurrent Systems: CSP

- Established formalism to describe concurrent systems.
- Still research on foundations; applications in industry, e.g. Train Controllers, Avionics, Security Protocols.
   Roscoe. The Theory and Practice of Concurrency. Prentice Hall, 1998.

Abdallah, Jones, Sanders (eds). CSP: The First 25 Years. Springer 2005.

### Modelling Data: CASL

- CASL = <u>Common Algebraic Specification Language</u>.
- De-facto standard in algebraic specification Mosses (ed). CASL Reference Manual, Springer 2004. Bidoit, Mosses. CASL User Manual, Springer 2004.

### A C++ template

```
Template:
```

template <typename T>
T square(T x) { return x \* x; }

Instantiation:

```
square <int>
```

Checks:

- Extract the signature required for T: "type T, \* ".
- Check that <int> offers this signature.

# The same in the specification language CASL

#### Generic specification:

spec MYTEMPLATE [sort T op \_\_\*\_:  $T \times T \to T$ ] = op square :  $T \times T \to T$  $\bullet \forall x, y : T \bullet square(x, y) = x * y$ 

**Instantiation:** MYTEMPLATE [INT] Checks:

INT is a refinement of
sort T op \_\_\_\*\_: T × T → T

(in our example: boils down to a check on signatures only)

### **Underlying Framework: Institutions**

Goguen, Burstall. Institutions: Abstract model theory for specification and programming. 1992.

#### Institutions speak about

- Signatures (e.g.: T is a type, \* is an operation)
- Models (e.g.: interpretation of type T by set **Z**)
- Formulae (e.g.: square(x, y) = x \* y
- Satisfaction (e.g.:  $\mathbf{Z} \models x * y = y * x$ )

### **C++ Concepts with threads?**

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#### Generic specification:

spec MyTemplate
 [call1 -> call2 -> SKIP [] call2 -> call1 -> SKIP [T= P]
= P; call3 -> Skip

Instantiation: MyTemplate [call1 -> SKIP ||| call2 -> SKIP]

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#### Check:

call1 -> SKIP ||| call2 -> SKIP
 is a refinement of
 call1 -> call2 -> SKIP [] call2 -> call1 -> SKIP

### Questions

- $\bullet$  What properties of threads make sense for C++ Concepts?
- How do we formulate properties of threads?
- $\bullet$  What is a useful "refinement" on C++ threads?

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### Suggestion:

A "process" algebra of C++ threads - formulated as an institution.

## A testing scenario

#### Tests on various levels:

Concept (abstract tests)	
refinement	
	ancement
Instantiated Concept	System using the Instantiated concept
(concrete tests)	(concrete tests in a context)

#### Refinement and Tests:



### Links to publications

- Mossakowski, Roggenbach: *Structured CSP A Process Algebra as an Institution*. 2007.
- Mossakowski, Roggenbach: An institution for processes and data. 2008.
- Kahsai, Roggenbach, Schlingloff: Specification-based testing for refinement. 2007.
- Kahsai, Roggenbach, Schlingloff: Specification-based testing for Software Product Lines. 2008.